

Lecture Notes

L- Lipschitz Continuity $\|f(x)-f(y)\| \leq L\|x-y\|$

L-smoothness $\| \nabla f(x) - \nabla f(y) \| \leq L \| x - y \|$ iff $-L \leq \nabla^2 f(x) \leq L$ iff $| \lambda | \leq L$

proof: $\| \nabla f(x) - \nabla f(y) \| = \| \nabla^2 f(\xi) (x-y) \| \leq L \| x-y \|$

Thm. f is L -smooth $\rightarrow f(y) \leq f(x) + \nabla f(x)^T (y-x) + \frac{L}{2} \| y-x \|^2$ (Taylor Expansion)

Thm. GD: $f(x_{k+1}) - f(x_k) \leq -\epsilon (1 - \frac{\epsilon}{2L}) \| \nabla f(x_k) \|^2 \leq -\frac{\epsilon}{2L} \| \nabla f(x_k) \|^2$ Significance: $\| \nabla f(x_k) \|^2 \leq \frac{2}{\epsilon} (f(x_k) - f(x^*))$

$$f(x_k) - f(x^*) \leq \frac{2L}{\epsilon} \| \nabla f(x_k) \|^2 \text{ using } \| \nabla f(x_k) \|^2 \leq \frac{2}{\epsilon} (f(x_k) - f(x^*))$$

m -strongly convex $f(x) - \frac{m}{2} \| x \|^2$ is convex iff $f(y) \geq f(x) + \nabla f(x)^T (y-x) + \frac{m}{2} \| y-x \|^2$ iff $\nabla^2 f(x) \geq mI$ iff $\lambda \geq m$

Thm. $\| x_{k+1} - x^* \|^2 \leq (1 - m\epsilon) \| x_k - x^* \|^2 \leq (1 - m\epsilon)^k \| x_0 - x^* \|^2$

Some Tuning

1. $x_k = \text{argmin}_x \{ f(x) - \frac{\epsilon}{2} \| x - x^* \|^2 \}$ $f(x_k) - f(x^*) \leq (1 - \frac{\epsilon}{2L})^k (f(x_0) - f(x^*))$ (2 Inequalities + Minimization)

2. Armijo's rule: backtracking line search $f(x_k) - f(x_k + \alpha d_k) \geq \alpha \epsilon \| \nabla f(x_k) \|^2 \| d_k \|^2$ (usually using $\alpha \leftarrow \beta \alpha$)

3. AGD $x_{k+1} = \frac{1}{2} (x_k + \nabla f(x_k))$, $x_{k+1} = \text{prox}_{\frac{1}{2L}}(x_k - \nabla f(x_k))$

4. Newton $x_{k+1} = x_k - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$

$$\text{Thm. } \| x_k - x^* \| \leq \frac{2L}{\epsilon} \left(\frac{2L}{\epsilon} \| x_0 - x^* \| \right)^{2^k} \quad \| A \|_{a,b} = \max_{\| x \|_a = 1} \| Ax \|_b$$

Damped Newton's method: $f(x_{k+1}) = f(x_k) + \epsilon \nabla f(x_k)^T d$

Proximal GD

$$x_{k+1} = \text{argmin}_x \frac{1}{2} \| x - (x_k - \nabla f(x_k)) \|^2 \rightarrow x_{k+1} = \text{argmin}_x \{ \frac{1}{2} \| x - (x_k - \nabla f(x_k)) \|^2 + h(x) \}$$

$\text{prox}_h(x) \triangleq \text{argmin}_x \{ \frac{1}{2} \| x - x_k \|^2 + h(x) \}$, then $x_{k+1} = \text{prox}_{h(x)}(x_k - \nabla f(x_k))$

Convergence Analysis $\| x^k - x^* \|_a \leq (1 - \frac{\epsilon}{2L})^k \| x^0 - x^* \|_a$

Lagrange (min f(x), Ax=b)

Thm. x^* is optimal iff $\nabla f(x^*) \perp \text{Null}(A)$

$$\text{Null}(A) = \{ x | Ax=0 \}, \text{Range}(A) = \{ Ax : x \in \mathbb{R}^n \}, \text{Null}(A)^\perp = \text{Range}(A^T)$$

$$L(x, \lambda) = f(x) + \lambda^T (Ax - b), \text{Lagrange Equation } \nabla L = 0$$

A Regular point x^* of a function h if $\nabla h(x^*) \neq 0$, \tilde{h} if $\nabla h(x^*)$ are linearly independent.

ICP: $g_j(x) \leq 0$

Karush-Kuhn-Tucker conditions

If x^* is a local minimum of ICP and also a regular point, then:

$$1. \nabla f(x^*) + \sum_{j=1}^m \lambda_j \nabla g_j(x^*) = 0$$

$$2. \lambda_j \geq 0$$

$$3. \lambda_j g_j(x^*) = 0 \text{ --- complementary slackness condition}$$

Sufficiency of KKT conditions for convex problem

Projected GD $x_{k+1} = \text{proj}_C(x_k - \nabla f(x_k)) \triangleq x_k - \beta \nabla f(x_k)$

Connection to proximal GD $I_C \triangleq \begin{cases} 0 & \text{if } x \in C \\ \infty & \text{if } x \notin C \end{cases}$, $\text{prox}_{I_C}(y) = \text{proj}_C(y)$

Newton's method for ECP: $\min_x f(x) = \frac{1}{2} x^T Q x + g^T x + c$ s.t. $Ax=b$ ($Q \succeq 0$, A has full column rank)

$$\text{KKT system } \begin{bmatrix} Q & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} x^* \\ \lambda^* \end{bmatrix} = \begin{bmatrix} -g \\ b \end{bmatrix}$$

$$\text{KKT matrix } \triangleq K \quad \text{Lemma: } \text{Null}(K) = \left\{ \begin{bmatrix} x \\ \lambda \end{bmatrix} : x \in \text{Null}(Q) \cap \text{Null}(A) \right\}$$

Thm. if KKT system has no solution, then QP is either infeasible or unbounded below

Thm. if K is nonsingular, there unique solution.

Following conditions are equivalent:

$$1. K \text{ is nonsingular} \quad 4. F^T Q F \succ 0 \text{ for any } F \in \mathbb{R}^{n \times (n-k)} \text{ s.t. } \text{Range}(F) = \text{Null}(A)$$

$$2. \text{Null}(Q) \cap \text{Null}(A) = \{0\} \quad \text{Particularly, if } Q \succ 0, \text{ then } K \text{ is nonsingular}$$

$$3. Ax=0 \wedge x \neq 0 \Rightarrow x^T Q x > 0$$

$$\min_d h(d) \triangleq f(x_k + d) - \nabla f(x_k)^T d + \frac{1}{2} d^T \nabla^2 f(x_k) d$$

$$\text{s.t. } A(x_k + d) = b \Leftrightarrow Ad = 0$$

$$\text{Now we have } \begin{bmatrix} \nabla^2 f(x_k) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} d \\ \lambda \end{bmatrix} = \begin{bmatrix} -\nabla f(x_k) \\ 0 \end{bmatrix}$$

Elimination $x = \tilde{x} + Fz$

Log barrier function $B(x) = -\sum_{i=1}^m \log(-g_i(x))$

$$\Rightarrow f(x) + \frac{1}{\epsilon} B(x) \quad \text{control point and control path } x^*(t), t \rightarrow \infty$$

Seems $\frac{1}{\epsilon}$ -suboptimal

solution: an increasing sequence ϵ_k with $x^*(\epsilon_k)$ as the initial point

Dual LP

$$\min_x f(x) = c^T x$$

$$\max_{\lambda, \mu} \psi(\lambda, \mu) = \lambda^T b + \mu^T h$$

$$\text{s.t. } Ax = b \quad \leftarrow \text{Dual}$$

$$\text{s.t. } A^T \lambda + G^T \mu = 0$$

$$Gx \geq h$$

$$\mu \geq 0$$

Weak duality $f(x) \geq \psi(\lambda, \mu)$

Strong duality $f(x^*) = \psi(\lambda^*, \mu^*)$

$$W_1(a, b) = \max_{\pi \text{ is a coupling}} E_{Y \sim \pi(Y)} E_{X \sim \pi(X)}$$

Wasserstein distance

Lagrange dual function

$$\phi(\tilde{x}, \tilde{a}) = \inf_{x \in D} (f(x) + \sum \lambda h(x) + \sum \mu g(x)) \quad \text{downplay the role of } D$$

Always concave. (λ, μ) is feasible if $\lambda \geq 0$ and $\phi(\lambda, \mu) > -\infty$

Slater's condition

convex problem:

$$\min_x f(x)$$

$$\text{s.t. } g_j(x) \leq 0, j=1, 2, \dots, m$$

$$Ax = b$$

The above problem is strictly feasible, i.e. $\exists x \in D$, s.t. $g_j(x) < 0 \wedge Ax = b$.

Refined: g_j is affine $\rightarrow g_j(x) \leq 0$

Thm. Strong duality holds for (CP) under Slater's condition.

KKT conditions

KKT conditions for convex problems

Consider a differentiable convex problem and its dual,

$$\begin{array}{l|l} \min_x f(x) & \max_{\lambda, \mu} \phi(\lambda, \mu) = \inf \mathcal{L}(x, \lambda, \mu) \\ \text{s.t. } g(x) \leq 0 & \text{s.t. } \mu \geq 0 \\ h(x) = Ax - b = 0 & \end{array}$$

Theorem. KKT conditions hold at x^* with Lagrange multipliers λ^*, μ^* ,

1. (primal feasibility) $h(x^*) = 0, g(x^*) \leq 0$

2. (dual feasibility) $\mu^* \geq 0$

3. (stationarity) $\nabla_x \mathcal{L}(x^*, \lambda^*, \mu^*) = 0$

4. (complementary slackness) $\mu_j^* g_j(x^*) = 0, j=1, 2, \dots, m$

if and only if all the following conditions hold,

1. strong duality holds, i.e. $f^* = \phi^*$

2. x^* is a primal optimal solution, i.e. $f^* = f(x^*)$

3. (λ^*, μ^*) is a dual optimal solution, i.e. $\phi^* = \phi(\lambda^*, \mu^*)$

For Review

Positive definite $A > 0$ iff $\forall \lambda > 0$ iff $\lambda_1 > 0$ iff $|\det(A)| > 0$

[Convexity] $\partial x + (1-\theta)y$

Affine set $\{x \mid Ax=b\}$ Hyperplane $a^T x = b$ Halfspace $a^T x \leq b$

Ellipsoids $\{x \mid (x-x_0)^T P^{-1}(x-x_0) \leq 1\} = \{x \mid Ax=b \wedge A^T x \leq b\}$ Polyhedron $\{x \mid Ax=b \wedge A^T x \leq b\}$

Convex combination $\sum \theta_i x_i$, where $\theta_i \geq 0, \sum \theta_i = 1$

Convex Hull $\text{conv } S = \{\sum \theta_i x_i \mid \theta_i \geq 0, \sum \theta_i = 1\}$

Convexity-preserving operations: Intersections, affine transformation

Def. int C, \bar{C} (closure), ∂C (boundary), $P_C(x) = \arg \min_{z \in C} \|x-z\|_2$

Thm. Supporting/Separating hyperplane (BW Thm., $C_1 \cap C_2 = \emptyset$)

[Convex functions] $f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$

Def. $\text{epi}(f) = \{(x,t) \mid t \geq f(x)\}$; f is m -strongly convex if $f: \mathbb{R}^n \rightarrow \mathbb{R}$ is convex

Criteria: Zeroth order condition $\forall x, y, g(x) = f(x+y)$ is convex, $D \cdot f(x) \leq f(x+Dy)$

First order condition $\forall x, y, f(x) \geq f(y) + \nabla f(y)^T (x-y)$

Second order condition $\forall x, \nabla^2 f(x) \succeq 0$ ($\nabla^2 f(x) \succ 0 \iff f$ is strictly convex)

Convexity preserving operations

$\sum \theta_i f_i$ | $f_1(x) \leq f_2(x) \implies f_1(x) + \dots + f_n(x)$, f_1 increasing, g is convex | max, f_1 or \min, f_2 | $\inf_{x \in S} f(x)$, convex, convex C

[Convex optimization] $\min f(x)$

s.t. $g(x) \leq 0, Ax=b \implies$ Constructing feasible set

Methods of Transformations: Eliminating / bounding slack variables

Optimality conditions $\nabla f(x^*) = 0, \nabla f(x^*)^T (x-x^*) \geq 0, \begin{cases} \nabla f(x^*) + \lambda^T \nabla g(x^*) = 0 \\ \lambda^T b = c \end{cases}$

LP

	Standard	Inequality
$\min_x c^T x$	$\min_x c^T x$	$\min_x c^T x$
s.t. $Bx \leq d$	s.t. $Ax = b$	s.t. $Ax \leq b$
$Ax = b$	$x \geq 0$	

$Ax \leq b \implies Ax = b, C \geq 0$

$x \implies x' - x'' \geq 0, x' \geq 0, x'' \geq 0$

$Ax = b \implies b \in \text{range } A$

regular point $x: \nabla g_1(x), \nabla g_2(x), \dots, \nabla g_m(x)$ linear independent

[Lagrange] $\mathcal{L}(x, \lambda, \mu) = f(x) + \sum \lambda_i h_i(x) + \sum \mu_j g_j(x) = f(x) + \lambda^T h(x) + \mu^T g(x)$ $\mathcal{L}(x, \lambda, \mu) \leq f(x)$ for $\mu \geq 0$

KKT conditions: 1. primal feasibility $h(x^*) = 0, g(x^*) \leq 0$

2. dual feasibility $\mu^* \geq 0$

3. stationarity $\nabla_x \mathcal{L}(x^*, \lambda^*, \mu^*) = 0$

4. complementary slackness $\mu_j^* g_j(x^*) = 0$

Duality: $\phi(\lambda, \mu) = \inf_{x \in D} \mathcal{L}(x, \lambda, \mu)$ Domain: $\{(\lambda, \mu) \mid \phi(\lambda, \mu) > -\infty\}$

Dual problem $\max_{\lambda, \mu} \phi(\lambda, \mu)$ s.t. $\mu \geq 0$

Strong duality $\phi^* = f^*$ (holds for LP)

Slater's condition ($\exists x \in \text{int } D$ s.t. $g_j(x) < 0$ for all nonlinear inequality constraints g_j) \implies strong duality

Thm. KKT \iff Strong duality + primal optimality + dual optimality

[Algorithms]

— smooth f

descent method $x_{k+1} = x_k + \alpha d_k$

descent direction $\begin{cases} d_k = -\nabla f(x_k) \\ d_k = -[\nabla^2 f(x_k)]^{-1} \nabla f(x_k) \end{cases}$

stepsize α_k $\begin{cases} \text{constant } (\frac{1}{L} \text{ for } L\text{-smooth}) \\ \text{exact line search } \implies \alpha_k = \arg \min_{\alpha} f(x_k + \alpha d_k) \\ \text{backtracking line search } \implies \alpha_k \leftarrow \alpha_{k-1} \text{ until } f(x_k + \alpha d_k) > f(x_k) + \alpha \epsilon \nabla f(x_k)^T d_k \end{cases}$

— constrained problem

elimination

Newton's method $\implies \begin{bmatrix} \nabla f(x) & \lambda^T \\ A & 0 \end{bmatrix} \begin{bmatrix} d \\ \lambda \end{bmatrix} = \begin{bmatrix} -\nabla f(x) \\ 0 \end{bmatrix}$ KKT system